The interrupt is received by the kernel, and in order to process it further interrupts are disabled. The current program is saved by the kernel and the attention shifts over to the interrupt the source of the interrupt is identified using the by the Interrupt Service routine (ISR) using the interrupt dispatch table (IDT).

The interrupt is executed and the program that was running previously is resumed and the save that was made at the start is used. The Interrupts are re-enabled again in case something ends up happening again.